## Expressing meaningful processing requirements among $HeTe_RoG^E$ neOu& nodes in an active network

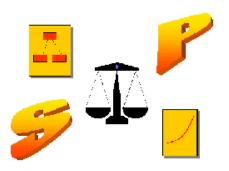
Virginie Galtier, Kevin Mills, Yannick Carlinet, Stefan Leigh, Andrew Rukhin

National Institute of Standards and Technology

http://w3.antd.nist.gov/active-nets







## Outline of the presentation

#### Problem:

- Context: What are active nets? What are they for?
- Why is it interesting to know the CPU resources requirement of an active application (AA)?
- What are the sources of variability in the execution time of an AA?
- Proposed solution:
  - Two models to characterize the processing requirements of an application on any active node
  - A mechanism to scale the models from one node to a different one
- Discussion and future work

## Active networks overview

- Active packets carry not only data but also the code to process them which is executed at active nodes.
- Example: an application that sends MPEG packets can specify an intelligent dropping algorithm to be applied at intermediate nodes if congestion is detected.
- Advantage: fast and easy deployment of customized network services.

# Why is it important to know the CPU resource requirements of an active application?

- Implication: in an active net the processing requirements can vary a lot from packet to packet.
- Without modeling, prediction, measurement and control, 3 threats:
  - a packet may consume excessive CPU time at a node, causing the node to deny services to other packets,
  - an active node may be unable to schedule its resources to meet the performance requirements of packets,
  - an active packet may be unable to select a path
     that can meet its performance requirements.

## Existing control solutions

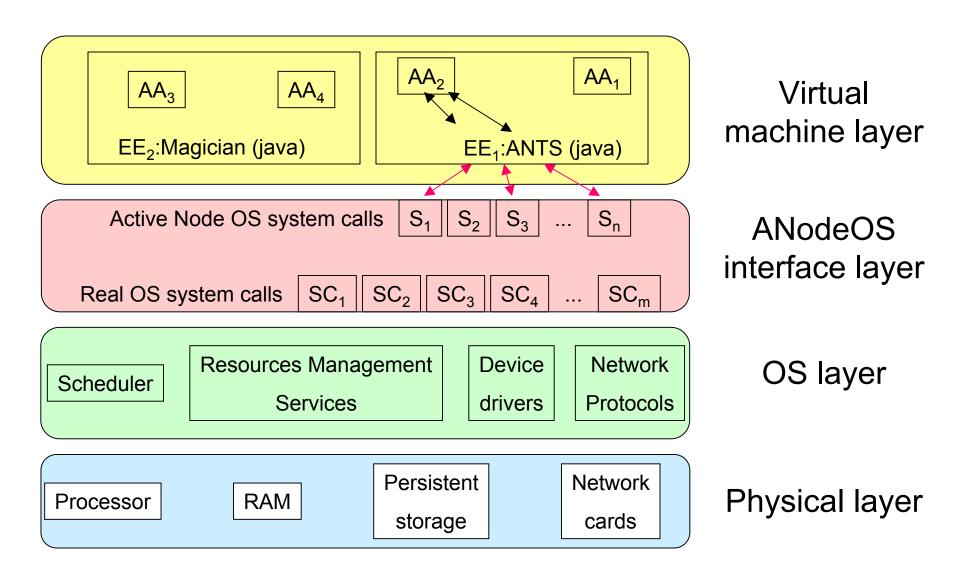
- A limit fixed by each node, the same for all packets.
- A time-to-live for the packet fixed by the application, the same for all nodes.
- Limitations with these solutions:
  - How to choose the limit?
  - This avoids major problems but doesn't permit optimum management.
  - Because: all applications are treated the same way.

## Necessity of modeling CPU requirements

- Idea to overcome these limitations: measure the CPU requirements of a packet once, and have the packet transport this information along with its data and code.
- Problem: there is no unit to measure CPU requirement that can be understood by all active nodes.

 It's necessary to have a model which captures all sources of variability and which can be translated on every node into a meaningful measure.

## Sources of variability in processing time



## Modeling active applications: trace

# Active Node OS System calls Monitoring



#### **Execution trace**

series of CPU time stamped system calls and transitions

```
AA_2
     EE<sub>1</sub>:ANTS (java)
       write
                        kill
read
   ANodeOS interface
         OS layer
      Physical layer
```

```
begin, user (4 cc), read (20 cc), user (18 cc), write(56 cc), user (5 cc), end

begin, user (2 cc), read (21 cc), user (18 cc), kill (6 cc), user (8 cc), end

begin, user (2 cc), read (15 cc), user (8 cc), kill (5 cc), user (9 cc), end

begin, user (5 cc), read (20 cc), user (18 cc), write(53 cc), user (5 cc), end

begin, user (2 cc), read (18 cc), user (17 cc), kill (20 cc), user (8 cc), end

...
```

## Modeling active applications: model M1

#### **Execution trace**



#### Model M1

(suited for ANTS applications)

```
begin, user (4 cc), read (20 cc), user (18 cc), write(56 cc), user (5 cc), end
```

begin, user (2 cc), read (21 cc),
user (18 cc), kill (6 cc), user
(8 cc), end

begin, user (2 cc), read (15 cc),
user (8 cc), kill (5 cc), user (9
cc), end

begin, user (5 cc), read (20 cc),
user (18 cc), write(53 cc), user
(5 cc), end

begin, user (2 cc), read (18 cc),
user (17 cc), kill (20 cc), user
(8 cc), end

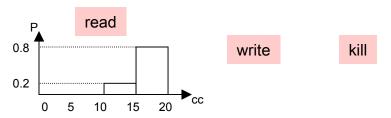
Scenario A:

sequence = "read-write", probability = 2/5

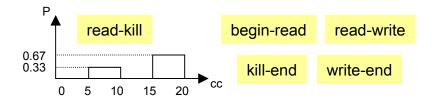
Scenario B:

sequence = "read-kill", probability = 3/5

Distributions of CPU time in system calls:



Distributions of CPU time between system calls:



## Modeling active applications: model M2

#### **Execution trace**



#### Model M2

(suited for Magician applications)

```
begin, user (4 cc), read (20 cc),
user (18 cc), write(56 cc), user
(5 cc), end

begin, user (2 cc), read (21 cc),
user (18 cc), kill (6 cc), user
(8 cc), end

begin, user (2 cc), read (15 cc),
user (8 cc), kill (5 cc), user (9
cc), end

begin, user (5 cc), read (20 cc),
user (18 cc), write(53 cc), user
(5 cc), end

begin, user (2 cc), read (18 cc),
user (17 cc), kill (20 cc), user
(8 cc), end
```

#### Scenario A:

```
Sequence = begin, user (4,5 cc), read (20 cc), user (18 cc), write (54,5 cc), user (5 cc), end probability = 2/5
```

#### Scenario B:

```
Sequence = begin, user (2 cc), read (18 cc), user (14.33 cc), kill (10.33 cc), user (8.33 cc), end probability = 3/5
```

## **Predicting CPU requirements**

 A node needs to predict not only the average CPU time required to execute a packet but also the high percentiles (example : 95% of executions are expected to complete within 70 cc).

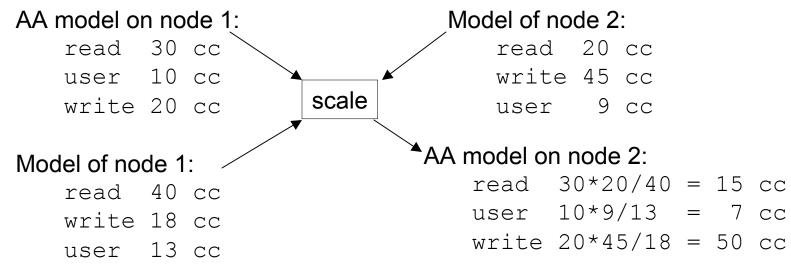
Model M1: simulation

Model M2: analytical computation

Active	Active	average absolute deviation of predictions from reality (%)										
Network	Application	M1, 100 bins, 20000 rep		M1, 50 bins, 20000 rep		M1, 50 bins, 500 rep		M2				
Platform		mean	high perc.	mean	high perc.	mean	high perc.	mean	high perc.			
ANTS	ping	0.859	0.9	0.643	1.622	2.696	9.8	0.028	16			
	multicast	0.398	1.94	0.351	3.002	4.913	15.93	0.001	18			
magician	ping	0.296	49	0.193	43			0.006	18			
	route	0.991	20	0.211	19			0.001	23			

## Overcoming node heterogeneity: node model

- Node model:
  - a system benchmark program ≤ for each system call, average system
  - for each EE, a user benchmark program < average time spent in the EE between system calls



To scale: a reference node model known by all other active nodes

## Overcoming node heterogeneity: results

Platform	Application	node 1	node 2	mean	high perc.
		Daisy	Blue	2.78	4.91
	Ping	Daisy	Sloth	4.55	11.05
		Blue	Daisy	3.63	5.64
ANTS		Sloth	Blue	7.69	8.33
		Daisy	Blue	0.32	7.29
	Multicast	Blue	Daisy	3.15	11.79
		Sloth	Daisy	23.38	15.7
		Blue	Daisy	11.49	20.03
	Ping	Blue	Sloth	8.01	5.2
Magician		Daisy	Blue	7.3	37.92
		Blue	Daisy	2.23	19.23
	Route	Daisy	Blue	1.59	34.54
		Sloth	Blue	19.04	44.3

## Limitations of our models

- Models can be large: O(number of scenarios, number of bins, distributions of the times).
- Simulation can be resource and time consuming: O(number of repetitions, size of the model).
- Trace-based models might represent probabilities not met in reality, if the scenario mix used to generate the traces does not represent the scenario mix actually seen on the nodes.
- Application behavior, such as looping, may depend on conditions at network nodes, and these conditions can be difficult to predict when generating the original traces.

### **Future work**

- Increase the test bed size (more nodes, more platforms, more applications)
- Investigate new models (your ideas are welcome!)
  - e.g., parameterize paths for loops
- Investigate an "Active" model:
  - gains experience as it travels through the net,
  - continuously evaluate which of the available coexisting models or prediction systems is the most accurate to return the prediction.
- Integrate our models with GE network-resource prediction system.

## Your turn...



Questions, suggestions...

















